

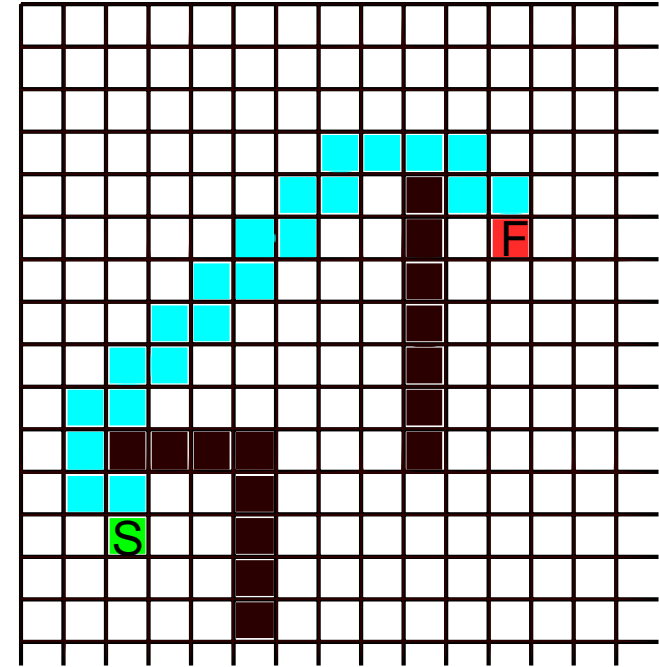
Solving Problems with Search

ITI0210, lecture 3 (2021)

Review

Pathfinding on 2D grid

- Expand BFS or best-first from start
- Solution is a sequence of steps
(up, left, up, ...)



<http://www.redblobgames.com/pathfinding/a-star/introduction.html>

Part I: Search space

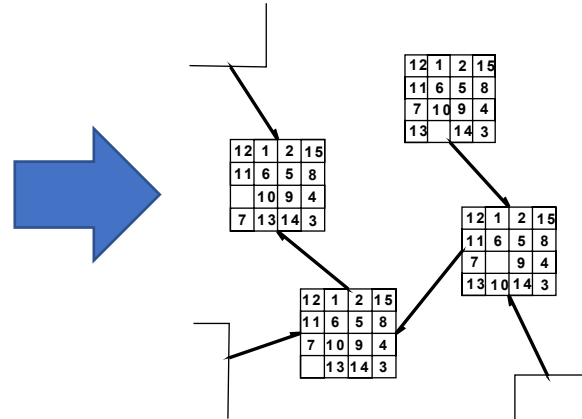
State space search

Actual problem



Physical

State space



Step 1: describe state space

Formulating a problem:

1. What are states? (grid squares)
2. State transitions? (moving left, right, up, down)
3. Where do I start? (square $x=2$, $y=3$)
4. What is the goal? (square $x=11$, $y=9$)

(blue: examples for 2D grid search)

Example formulation

The 8-Puzzle (smaller version of 15-puzzle)

Start and end states:

7	2	4
5		6
8	3	1

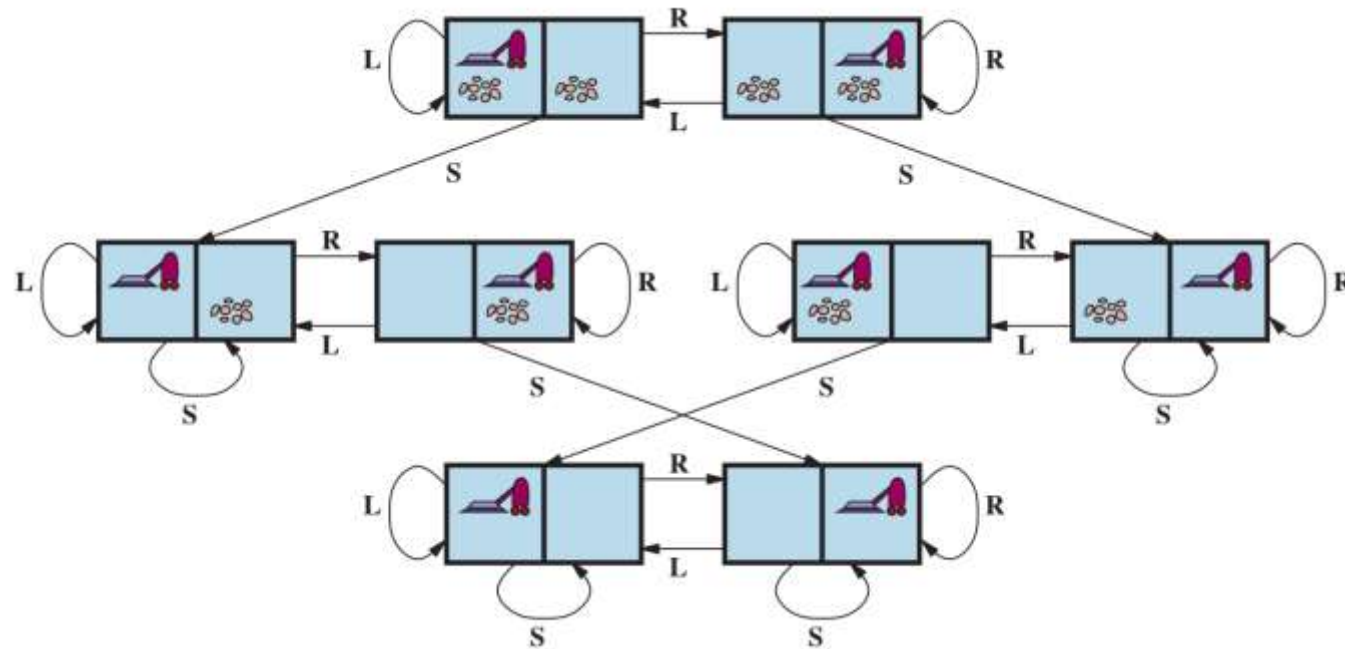
	1	2
3	4	5
6	7	8

State: one possible placement of pieces

State transitions: move a piece (e.g. “3 up” in start state)

Example formulation

Complete state space for the Vacuum World



L-left, R-right, S-suck (vacuum)

Step 2: implement in code

The 8-puzzle

State representation: array

map of array index to 3x3 coordinate:

0	1	2
3	4	5
6	7	8

Step 2: implement in code

The 8-puzzle

State representation: array

map of 3x3 coordinate to array index:

0	1	2
3	4	5
6	7	8

7	2	4
5		6
8	3	1



[7, 2, 4, 5, ~~0~~, 6, 8, 3, 1]

Step 2: implement in code

Move representation in 8-puzzle

What are the possible moves?

How to find them efficiently?

7	2	4
5		6
8	3	1

Step 2: implement in code

Move representation in 8-puzzle

Idea: “move” the empty spot

7	2	4
5		6
8	3	1

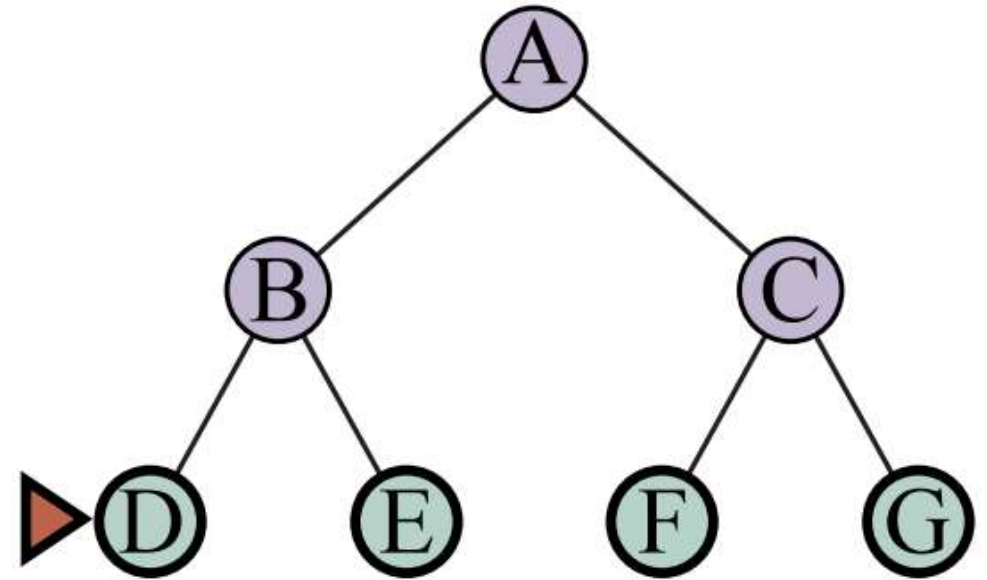
```
7
8 state = [7, 2, 4, 5, 0, 6, 8, 3, 1]
9
10 def actions(state):
11     empty = state.index(0)
12     # move up
13     up = empty - 3
14     if up >= 0:
15         # legal move, generate "UP" action
```

Exercise: what are $empty - 3$, $empty + 3$, $empty + 1$ etc?

Step 3: tree implementation

Node D needs:

- Linked state
- Path cost (A-B-D)
- Parent (B)
- Action (from B to D)

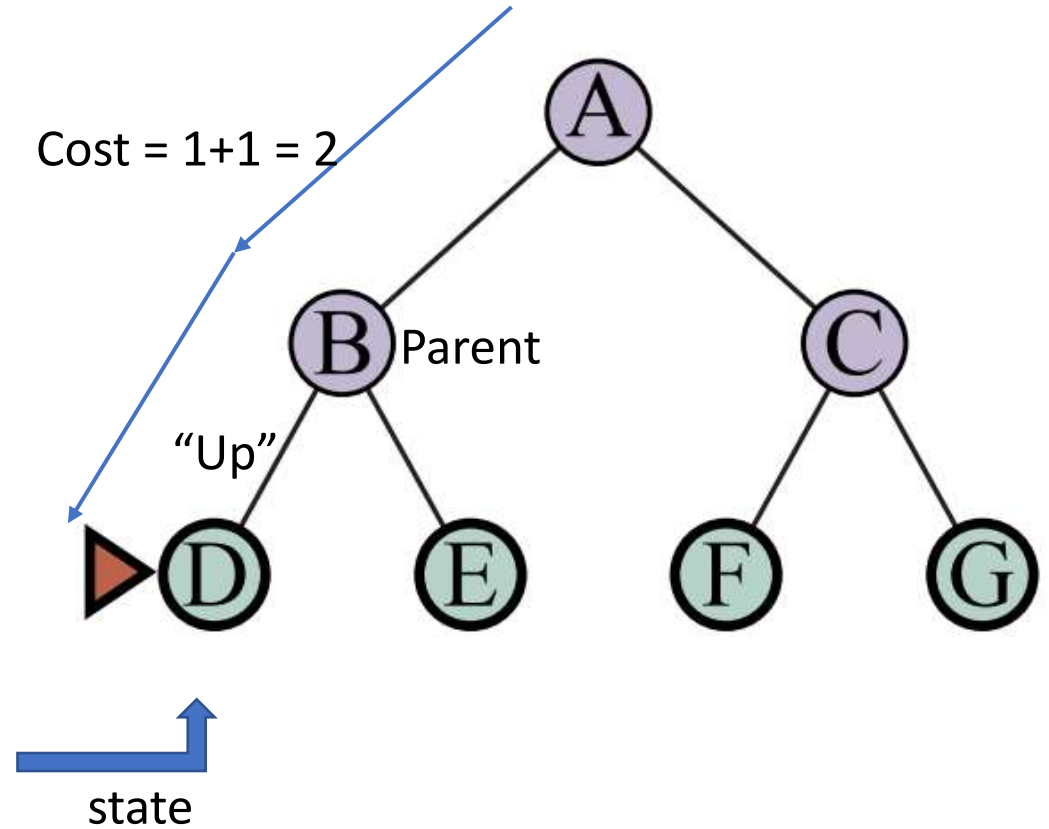


Step 3: tree implementation

Node D needs:

- Linked state
- Path cost (A-B-D)
- Parent (B)
- Action (from B to D)

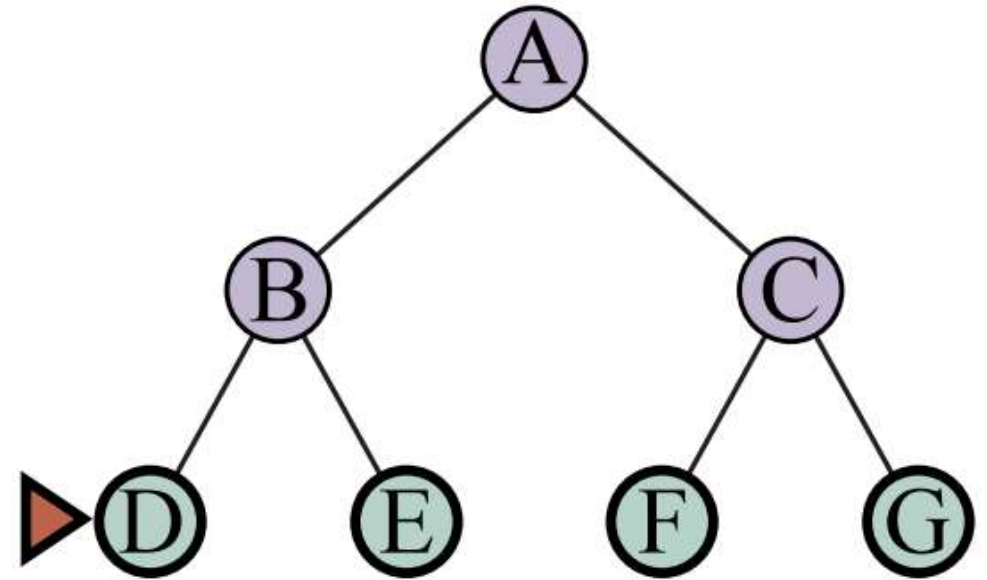
7	2	4
5		6
8	3	1



Step 3: tree implementation

Node D needs:

- Linked state
- Path cost (A-B-D)
- Parent (B)
- Action (from B to D)



Exercise: find these things from

<https://www.redblobgames.com/pathfinding/a-star/introduction.html>

Part II: Search algorithm

Generic search algorithm

SearchQueue() – keep nodes ordered by strategy

1	4	
start	node	5
3	6	

goal_check() – compare to goal state or condition

expand() – generate child nodes for a node

```
def search(start):  
    # start: node representing start state  
    queue = SearchQueue()  
    queue.add(start)  
    while not queue.empty():  
        node = queue.get()  
        if goal_check(node):  
            # solution node->parent->...->start  
            return node  
        for child in expand(node):  
            # insertion point depends on strategy  
            queue.add(child)  
    return "Failed"
```


Generic search algorithm

Breadth first, A* and other algorithms:

strategies of what direction to take in search

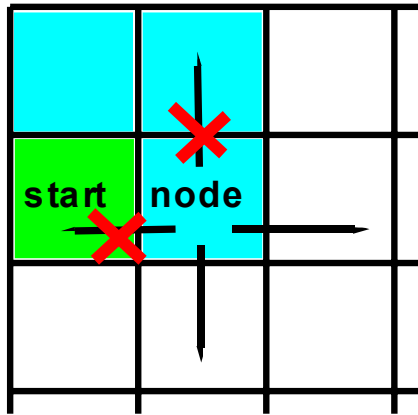
Queue order controls the search direction

```
def search(start):  
    # start: node representing start state  
    queue = SearchQueue()  
    queue.add(start)  
    while not queue.empty():  
        node = queue.get()  
        if goal_check(node):  
            # solution node->parent->...->start  
            return node  
        for child in expand(node):  
            # insertion point depends on strategy  
            queue.add(child)  
    return "Failed"
```


Generic search algorithm

Common optimization:

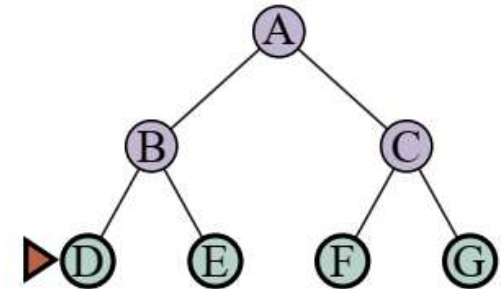
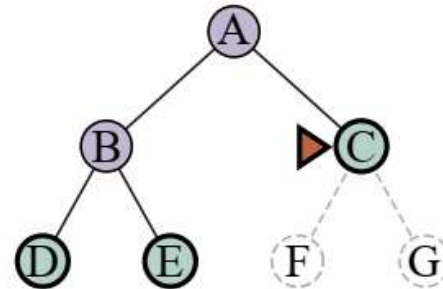
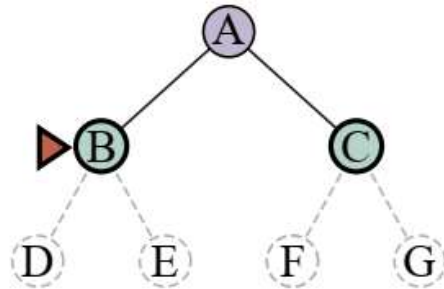
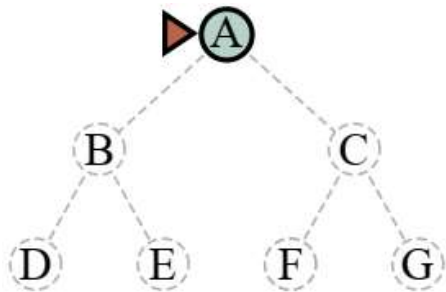
Keep track of visited states,
do not add them to queue



```
def search(start):  
    # start: node representing start state  
    queue = SearchQueue()  
    queue.add(start)  
    while not queue.empty():  
        node = queue.get()  
        if goal_check(node):  
            # solution node->parent->...->start  
            return node  
        for child in expand(node):  
            # insertion point depends on strategy  
            queue.add(child)  
    return "Failed"
```


Part III: Strategies

Breadth First Search



Queue:

<i>n</i>	A					
<i>d</i>	0					
<i>g</i>	0					

<i>n</i>	B	C				
<i>d</i>	1	1				
<i>g</i>	1	1				

<i>n</i>	C	D	E			
<i>d</i>	1	2	2			
<i>g</i>	1	2	2			

<i>n</i>	D	E	F	G		
<i>d</i>	2	2	2	2		
<i>g</i>	2	2	2	2		

Search step: remove 1st node and expand; new nodes to **end** of Q

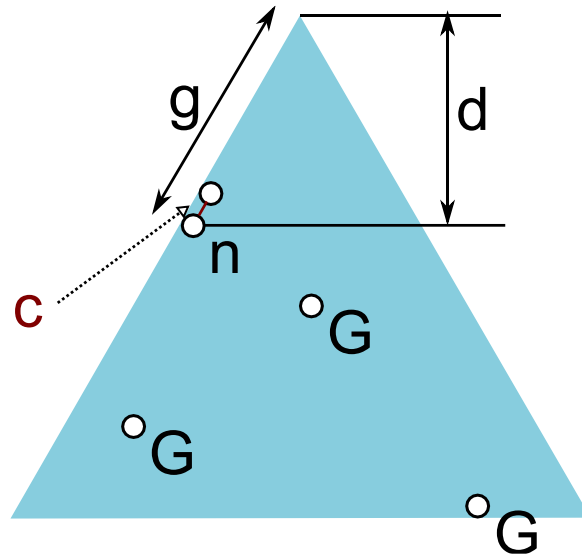
Search tree notation

n – node

d – node depth

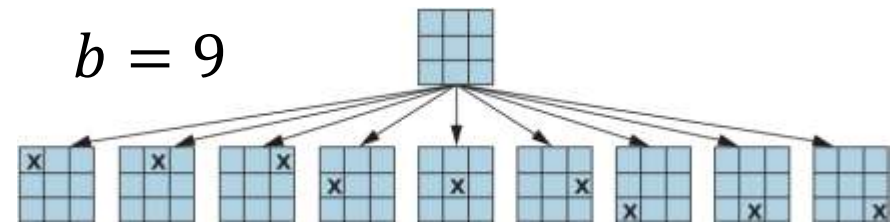
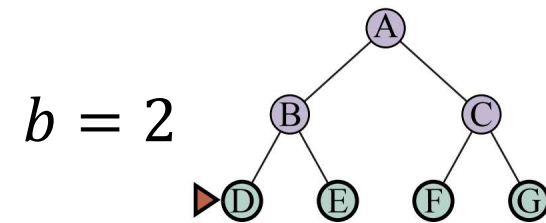
c – step cost

g – path cost



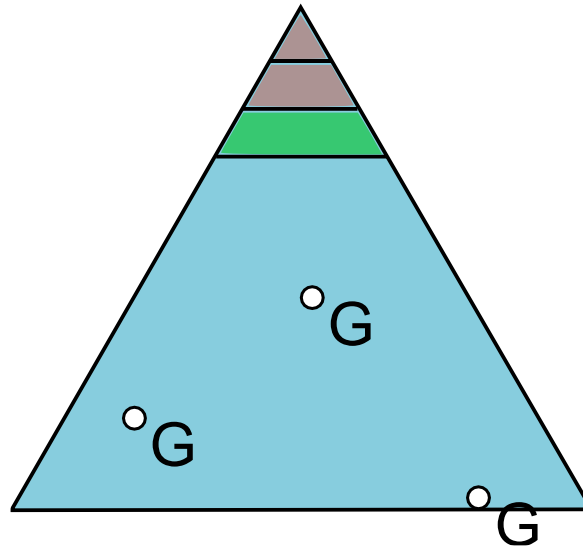
G – goal node(s)

b – branching factor



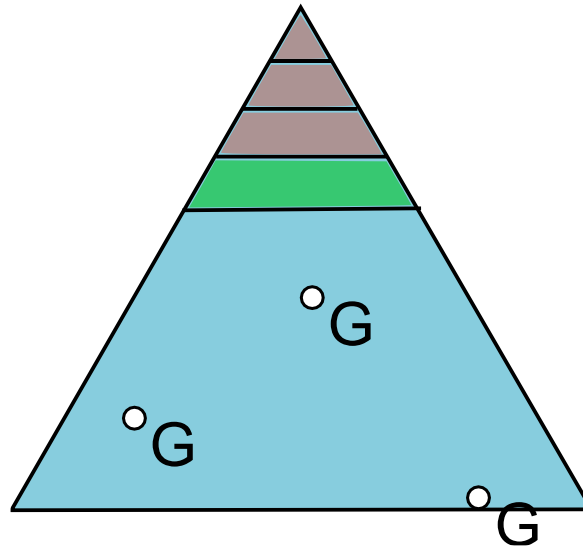
Breadth First Search

Progression



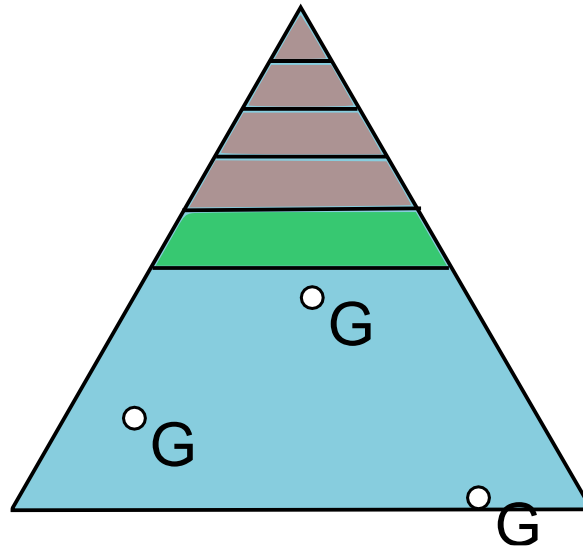
Breadth First Search

Progression



Breadth First Search

Progression

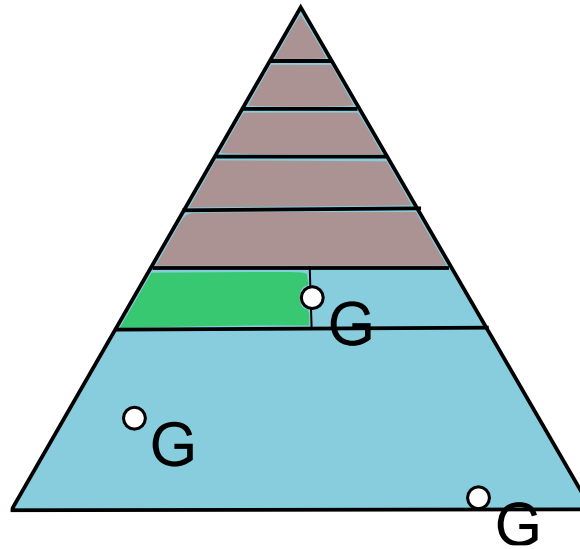


Breadth First Search

Goal found:

Minimal depth goal –
Optimal if we care about
number of steps

Always finds
finite depth goal

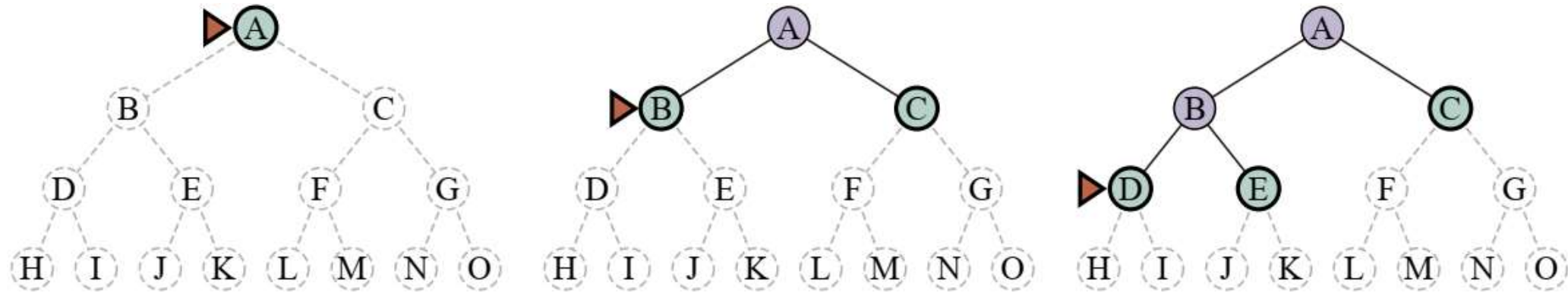


Space needed $O(b^d)$
(d – goal depth)

Time needed $O(b^d)$

(assuming goal check
before Q insertion)

Depth First Search



Queue:

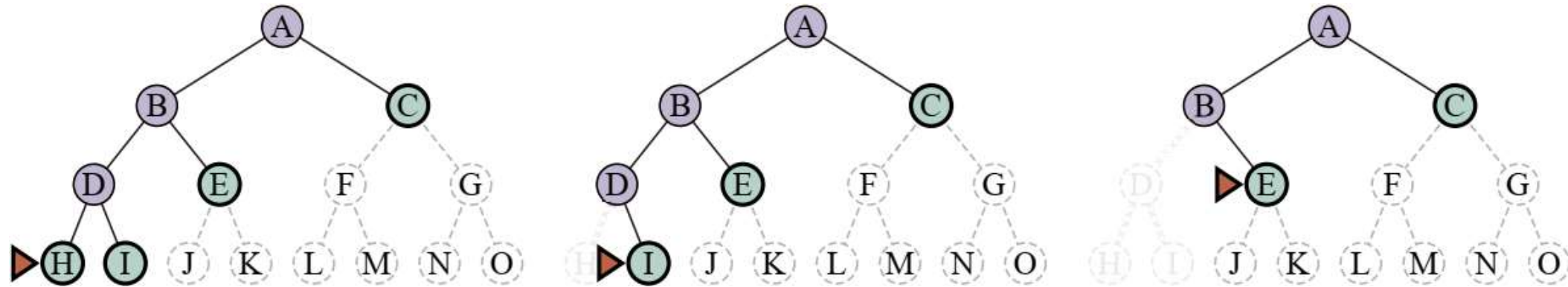
<i>n</i>	A					
<i>d</i>	0					
<i>g</i>	0					

<i>n</i>	B	C				
<i>d</i>	1	1				
<i>g</i>	1	1				

<i>n</i>	D	E	C			
<i>d</i>	2	2	1			
<i>g</i>	2	2	1			

Search step: remove 1st node and expand; new nodes to **start** of Q

Depth First Search



Queue:

<i>n</i>	H	I	E	C		
<i>d</i>	3	3	2	1		
<i>g</i>	3	3	2	1		

<i>n</i>	I	E	C			
<i>d</i>	3	2	1			
<i>g</i>	3	2	1			

<i>n</i>	E	C				
<i>d</i>	2	1				
<i>g</i>	2	1				

Search step: remove 1st node and expand; new nodes to **start** of Q

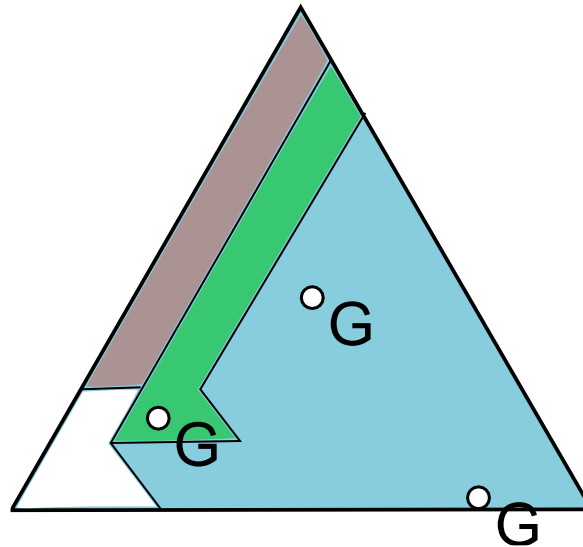
Depth First Search

Not Optimal

(may “shoot past”
optimal path)

Finite depth goal

not guaranteed (∞ tree)



Space needed $O(bm)$
(m – max depth)

Time needed $O(b^m)$

Depth First Search

How to handle ∞ trees?

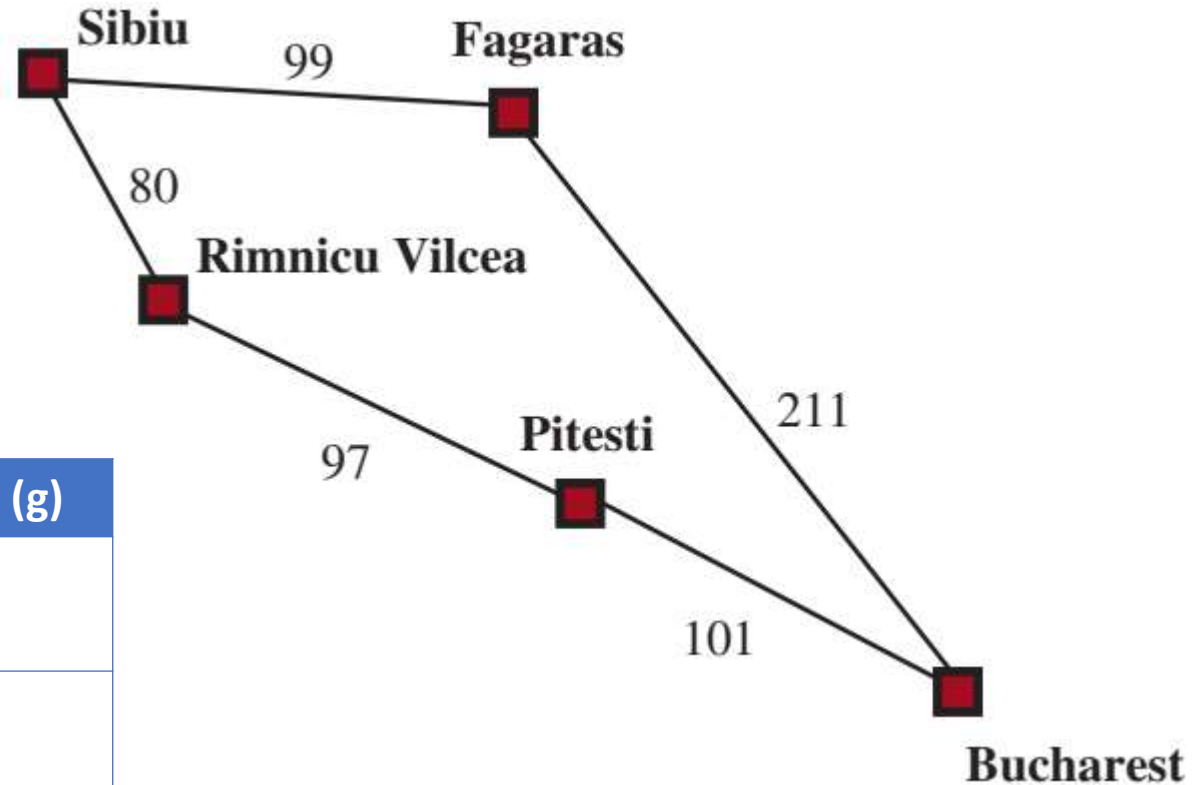
- Set depth limit
- Call DFS repeatedly with limit 1, 2, ...

This variant is called
Iterative Deepening Search (IDS)

Dijkstra Search

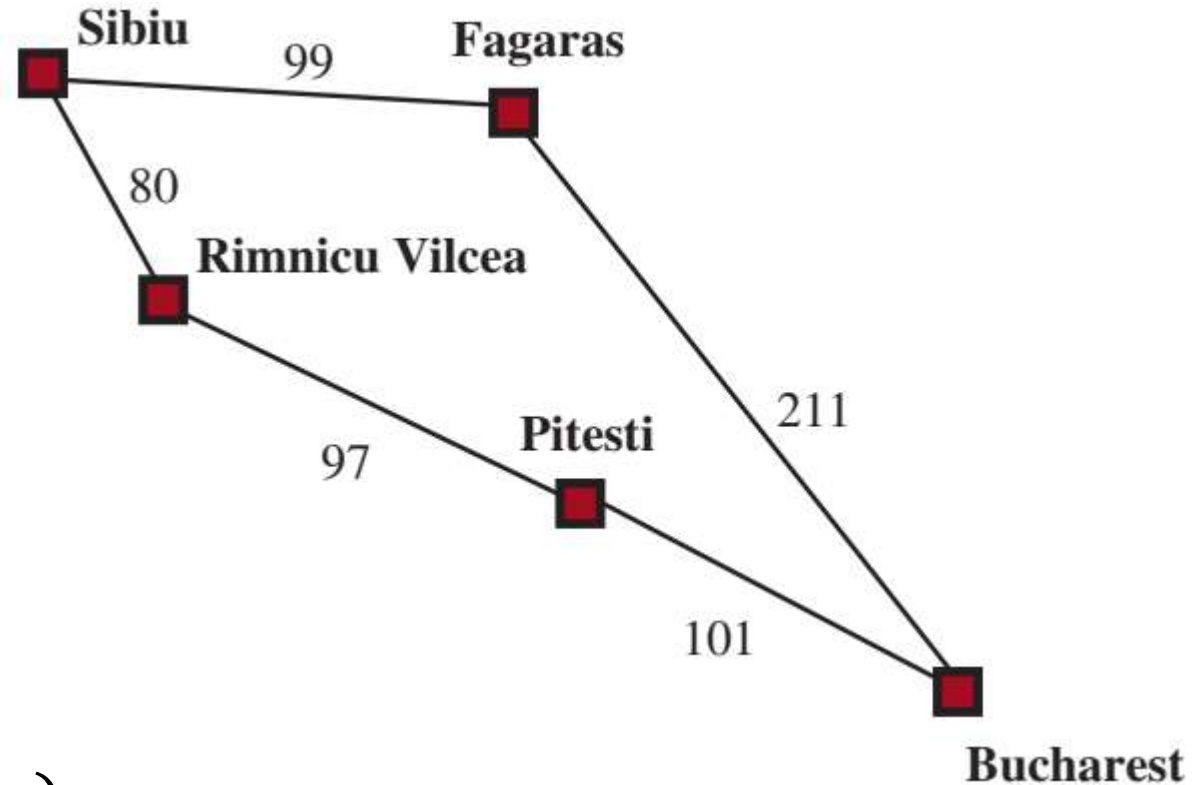
From “Romania” example
In AIMA
(find path to Bucharest)

Solution	depth	Path cost (g)
Sibiu, Faragas, Bucharest	2	310
Sibiu, R. V., Pitesti, Bucharest	3	278



Dijkstra Search

Node	d	g
Sibiu	0	0

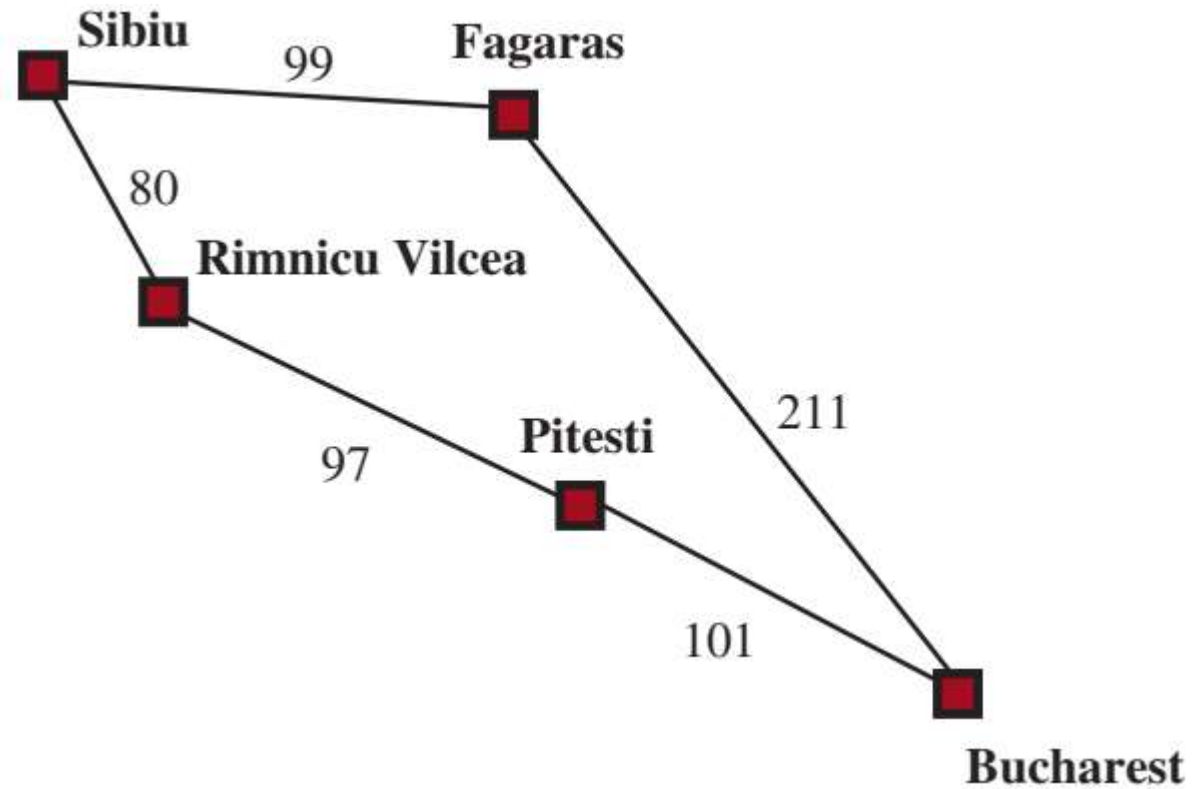


Search strategy:

Queue nodes in order of $g(n)$

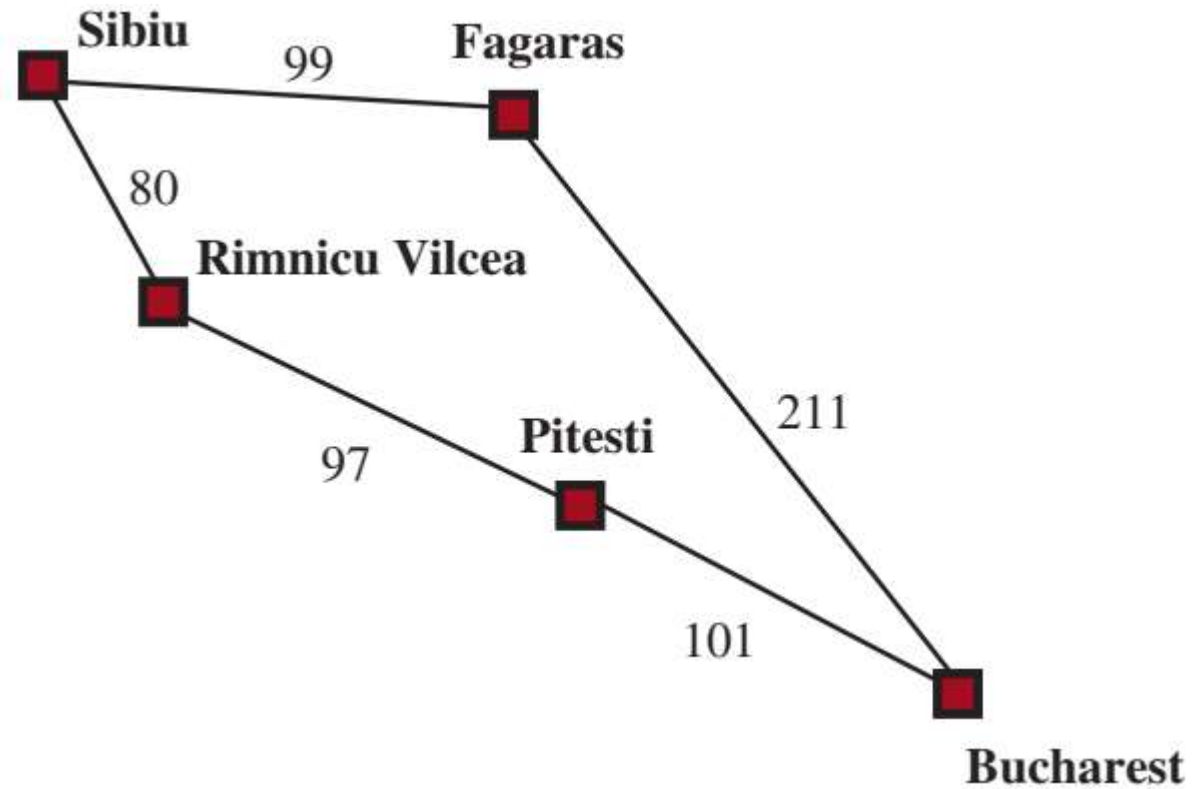
Dijkstra Search

Node	d	g
Rimnicu Vilcea	1	80
Fagaras	1	99



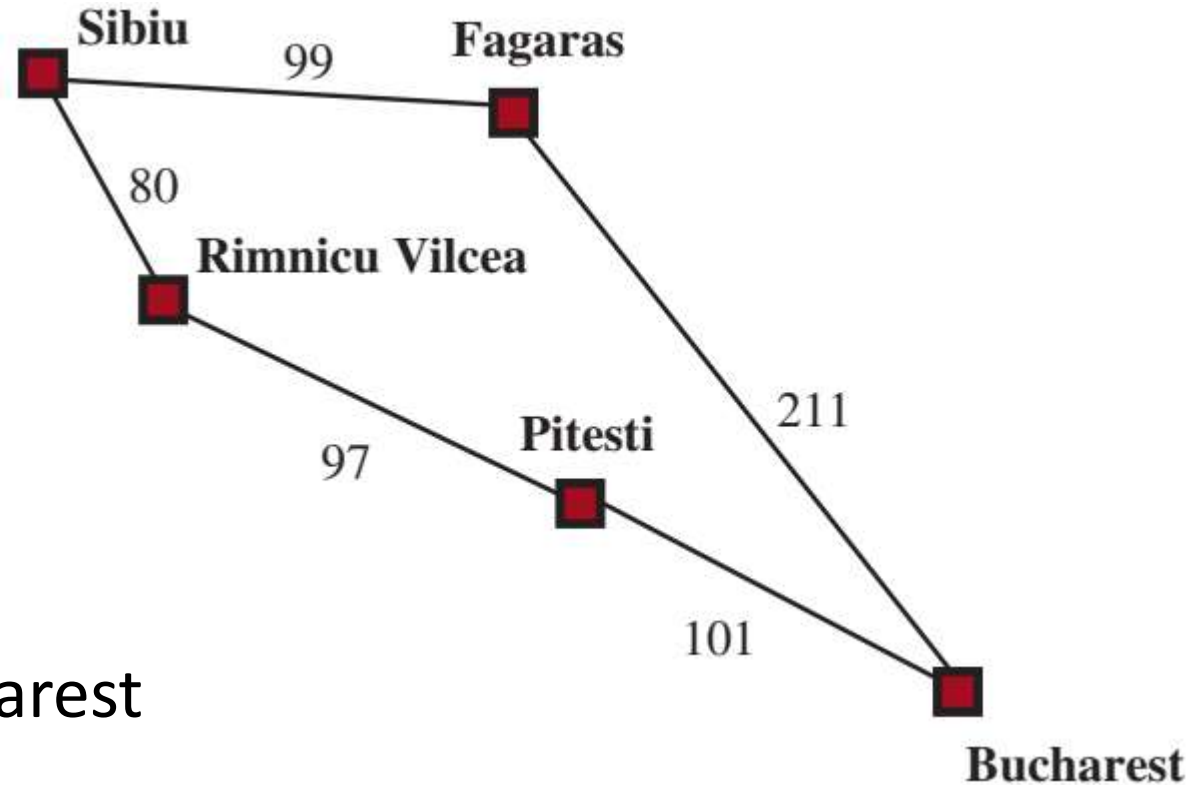
Dijkstra Search

Node	d	g
Fagaras	1	99
Pitesti	2	177



Dijkstra Search

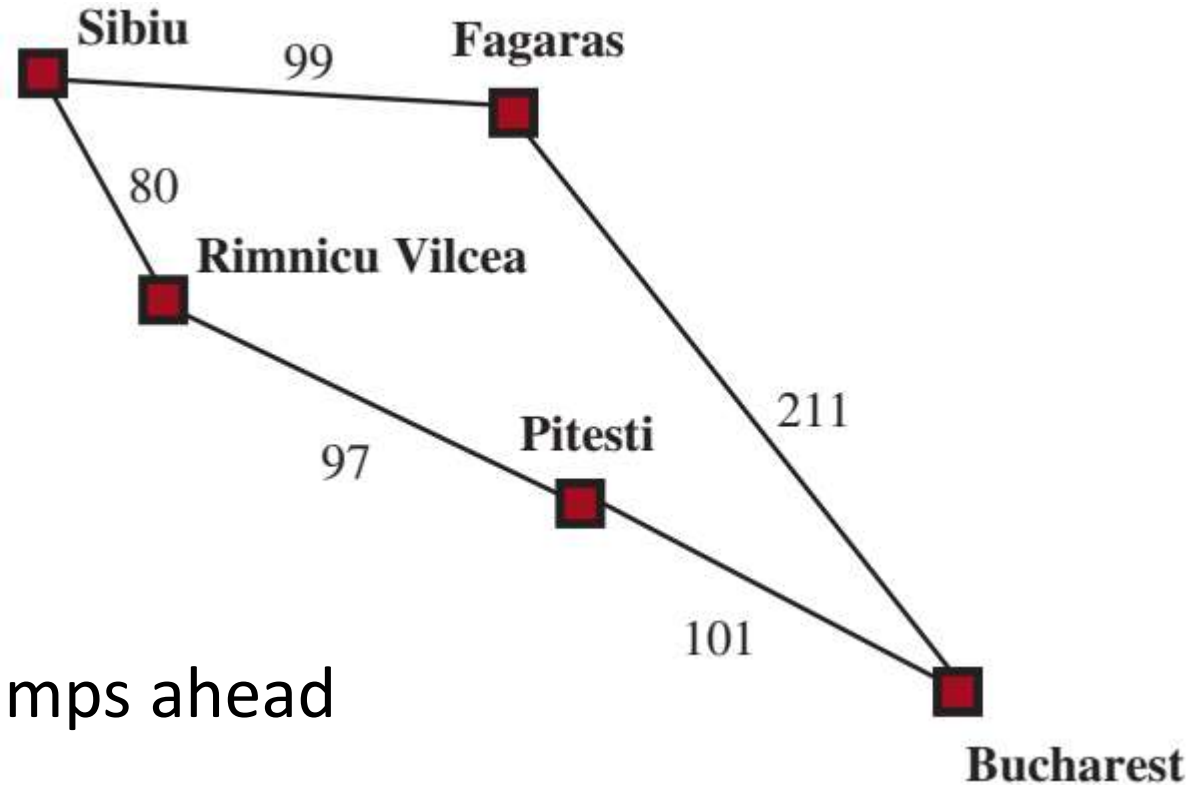
Node	d	g
Pitesti	2	177
Bucharest	2	310



Pitesti expands before Bucharest
(lower cost)

Dijkstra Search

Node	d	g
Bucharest	3	278
Bucharest	2	310

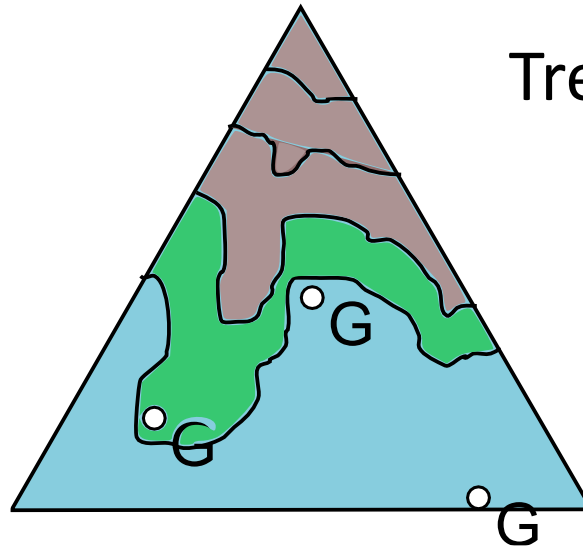


Second path to Bucharest jumps ahead
(lower cost)

Dijkstra Search

Optimal
(and uses cost!)

“Complete” -
Always finds
finite depth goal

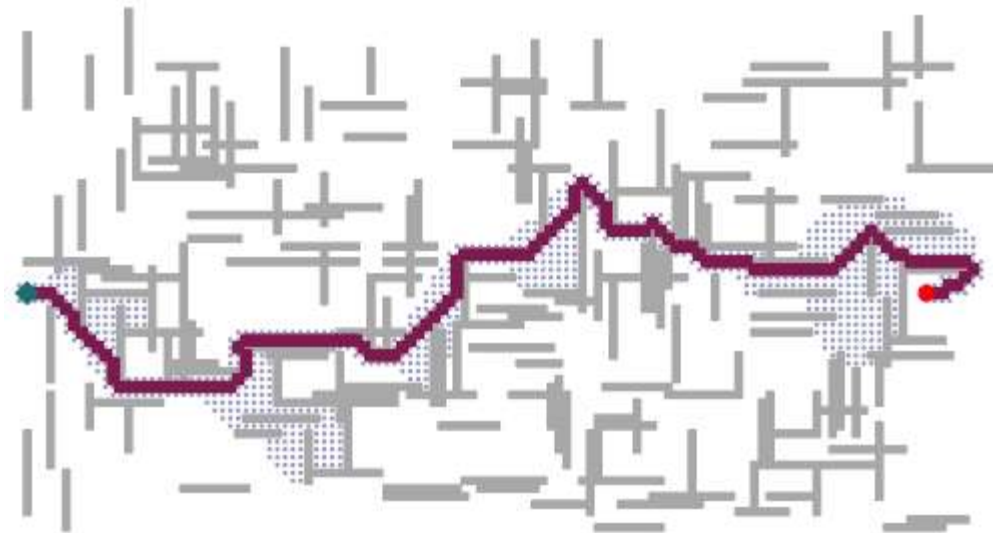


Tree is explored by cost $g(n)$
(not number of steps)

Heuristic search

If we know the direction of the goal
why search in the opposite direction?

Heuristic – some function to
measure closeness to goal



Heuristic Search

Order search queue by $f(n)$

Strategy	Queue order	Cares about
Dijkstra	$f(n) = g(n)$	Path cost
Greedy	$f(n) = h(n)$	Closeness to goal
A*	$f(n) = g(n) + h(n)$	Both!

$h(n)$ -distance from node to goal

Heuristic Search

Heuristic functions:

Recall: SLD, Manhattan for path finding on 2D maps

Exercise:

Find heuristics (how to estimate
of steps to goal) for this game:

(Solutions in AI/MA book 3.6)

7	2	4
5		6
8	3	1

	1	2
3	4	5
6	7	8

Heuristic search

$h(n)$ is “admissible” if it never **overestimates**

A* optimality proof:

<https://youtu.be/iTJvWfmp1vw?t=208>

Strategies summary

Strategy	Optimal for # of steps	Optimal for cost	Complete	Time	Space
BFS	Yes	No	Yes	$O(b^d)$	$O(b^d)$
DFS	No	No	No	$O(bm)$	$O(b^m)$
Dijkstra	Yes(1)	Yes	Yes		
Greedy	No	No	No		
A*	Yes(1)	Yes	Yes		

(1) Assuming cost doesn't matter, $c = 1$ achieves shortest # of steps